

# Demonstrating an Enhanced Ethernet Switch Supporting Video Sensing with Dynamic QoS

<http://www.ieeta.pt/lse/hartes>

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## Motivation

Video applications with **strict QoS requirements** are becoming increasingly common

- λ E.g. machine vision, object tracking, surveillance, driving aids
- λ These applications have several idiosyncrasies that must be taken into account:
  - Require a **large bandwidth**
  - Exhibit **high variability**, resulting from video compression
  - Have **dynamic behavior**
    - Sporadic activation (e.g. on-demand activation of video cameras in trucks and cars)
    - Acceptable performance can be attained with variable frame-rate, resolution, compression levels, ...
- λ Frequently demand **timely and predictable communications**

The **FTT-Enabled switch** features **hierarchical server-based traffic scheduling**, providing:

- λ **Disciplined use of the network**, thus coping with the large variability of compressed video sources;
- λ **Adaptability**, since allows changing the servers budget and thus varying dynamically the bandwidth allocated to each video stream
- λ **Reconfigurability**, since permits adding and removing servers dynamically, and consequently enabling and disabling video-streams

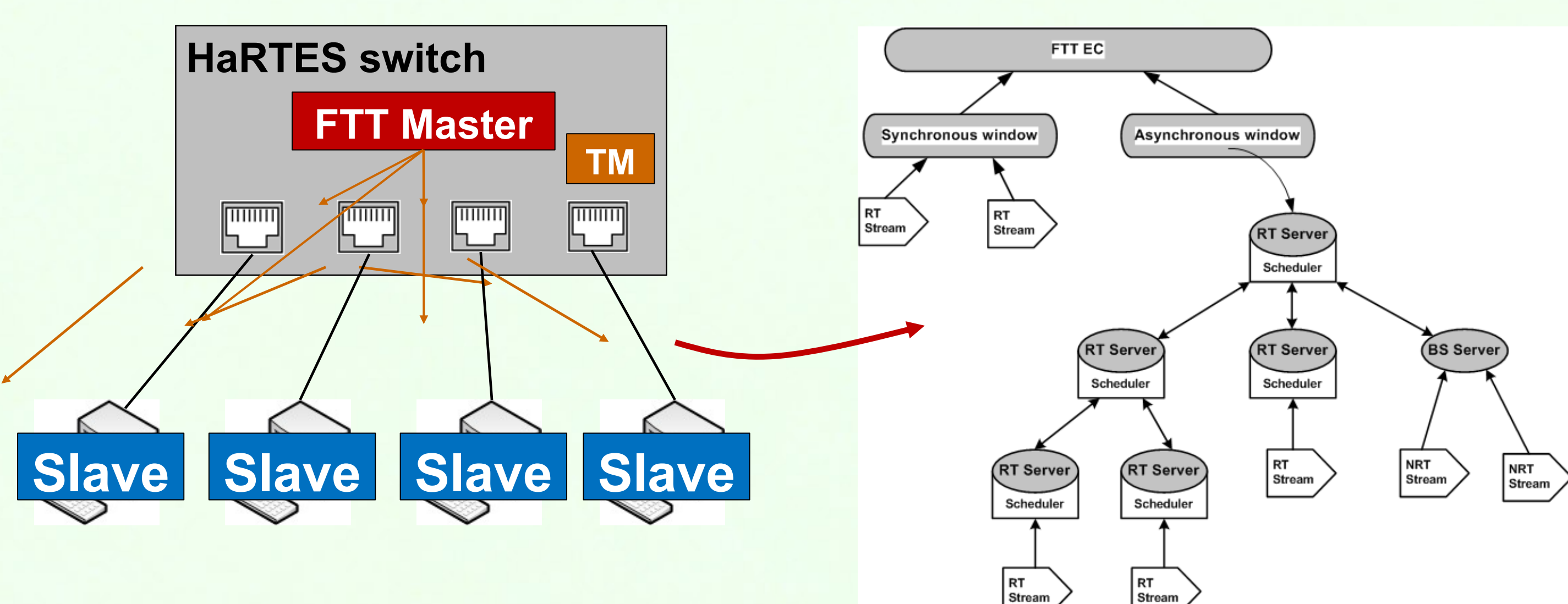
These features are demonstrated via a **simplified video surveillance application**

## Application Platform

The **server-based traffic scheduling** has been deployed using **HaRTES Switch** (FTT-enabled Ethernet Switch)

### HaRTES Switch:

- Based on FTT paradigm – master / multi-slave transmission control technique
- Communications organized in Elementary Cycles (divided in synchronous and asynchronous windows)
- The asynchronous window is handled by the presented **server framework**



## Traffic scheduling framework

A **flexible multi-level hierarchical server-based scheduling framework** for switched Ethernet that:

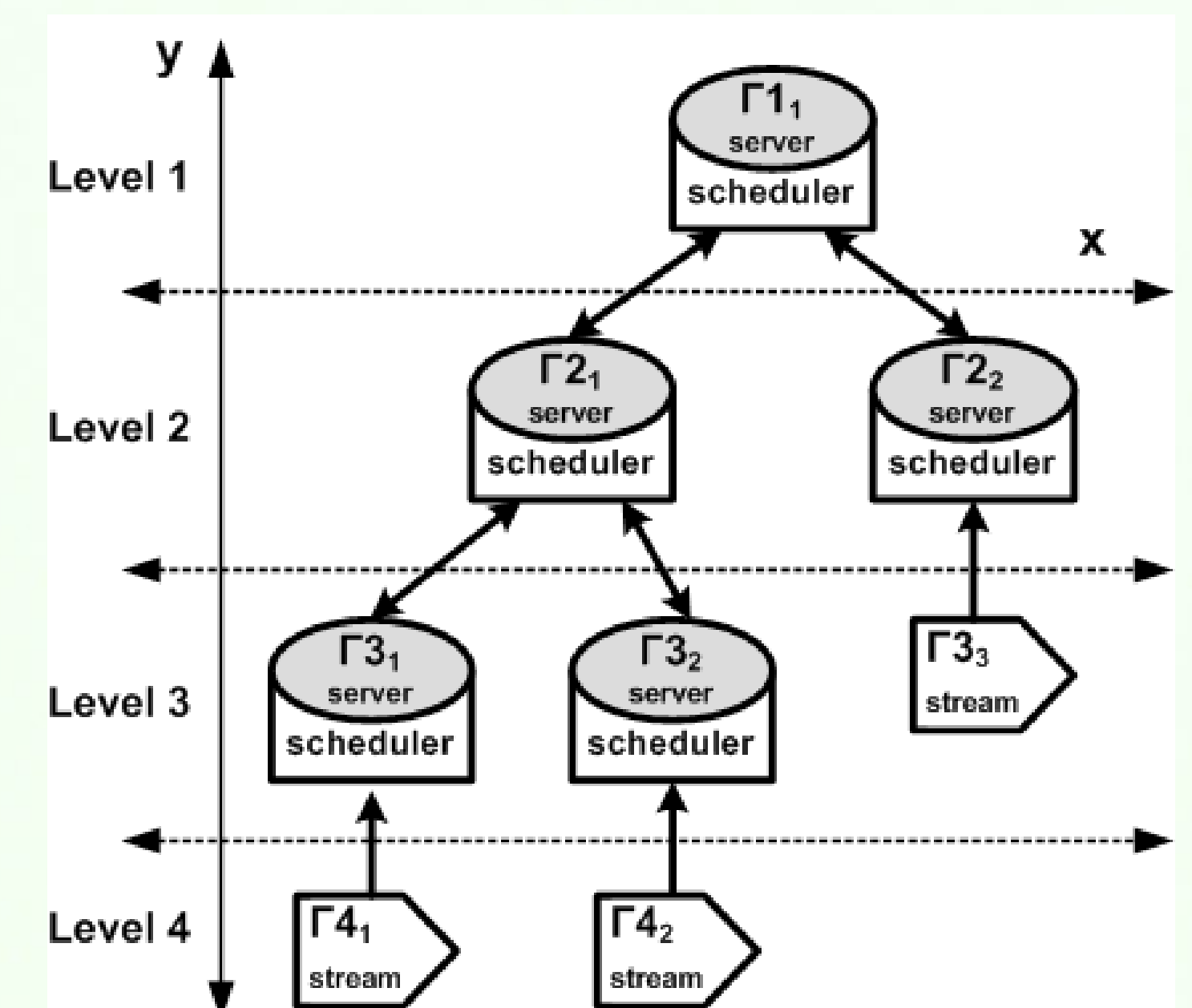
- Divides, subdivides and so on the network bandwidth in a hierarchical way, creating virtual channels
- Provides real-time guarantees and differentiated QoS (using admission control)
- Reconfigure and adapts the channel for varying flows
- Assures temporal isolation among channels

### Model Definition

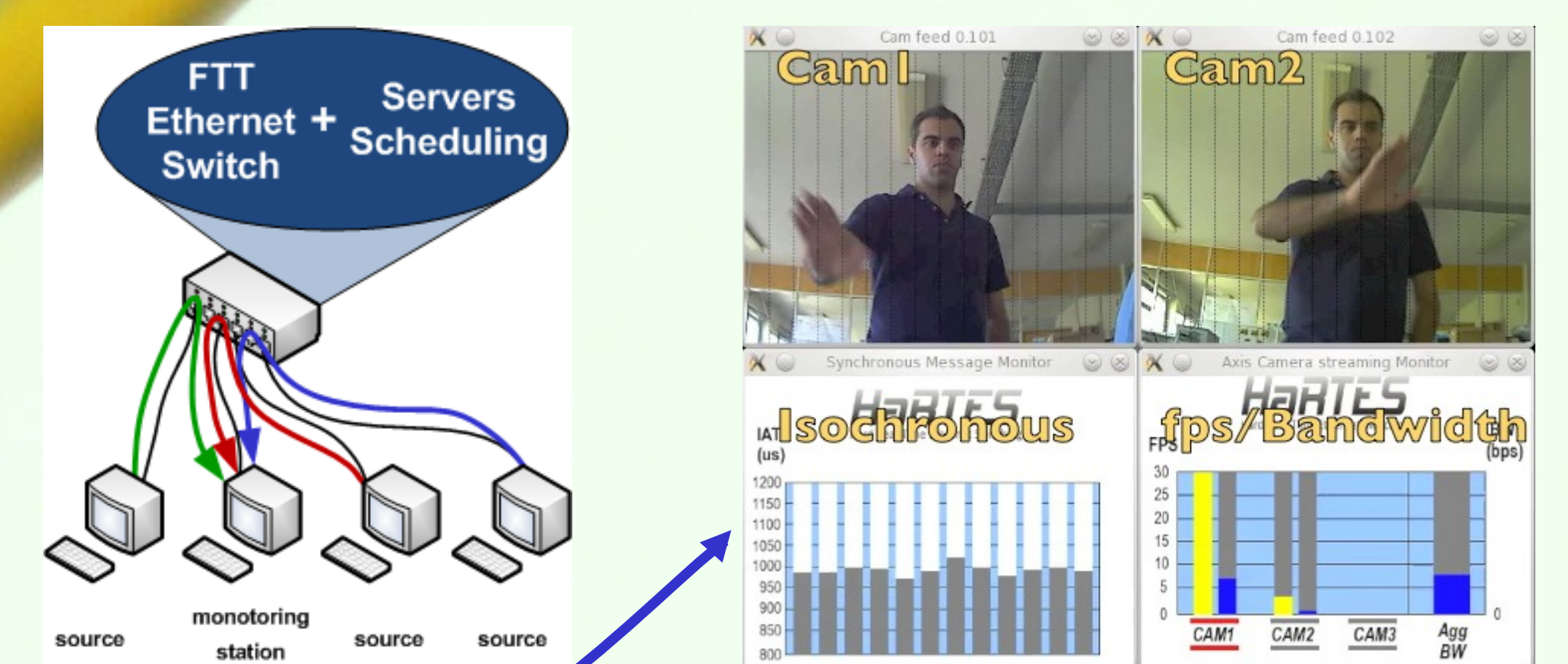
$$Srv_{y_x} = (C_{y_x}, T_{y_x}, Mmax_{y_x}, Mmin_{y_x}, P_{y_x}, RT_{y_x})$$

$$AS_{y_x} = (C_{y_x}, Tmit_{y_x}, Mmax_{y_x}, Mmin_{y_x}, P_{y_x}, RT_{y_x})$$

- $C_{y_x}$  - Maximum message transmission time
- $T_{y_x}$  - Period
- $Mmax_{y_x}$  - Transmission time - largest packet
- $Mmin_{y_x}$  - Transmission time - smallest packet
- $P_{y_x}$  - Parent component (server)
- $RT_{y_x}$  - Response time of each component



## Demonstration



**Demo 1** shows:

- The dynamic reconfiguration of the servers  
**Changes are made online without any impact on the streams not affected by the reconfiguration**
- Traffic isolation among the asynchronous streams  
**Sources are not allowed to use more bandwidth than the one granted by the server**

For that:

- 3 asynchronous video streams transmitted to the monitoring station

**Demo 2** shows:

- λ Strict periodicity of the isochronous services
- λ Traffic isolation between isochronous and asynchronous services  
**The bandwidth, latency and jitter of the isochronous traffic is kept constant, independently of the video streams load**

For that:

- 2 asynchronous video streams and one periodic data source concurrently transmitted to the monitoring station